# PasswordManager with CipherKey 

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#### Abstract

Transforming a plain text into non-readable format to keep a confidentially \& probity of data is called Encryption. And the method used to decode that into decipherable (readable) format, is called Decoding. To encode \& decode, algorithms were developed. This entire theory, entire method is called Cryptography. Many algorithms were developed, many were cracked, and many of the algorithms are still running nowadays also. So, here I came up with the new algorithm, with new method, with new idea in algorithm. Password managers are censorious (critical) pieces of software relied upon by users to securely store valuable and sensitive details from on-line banking passwords. Password manager taking action like a binary (digital) safe, which securely supply user's usernames, passwords and other sensitive information. The usernames and passwords are encoded before saving them into the database using cryptography methods.


Keywords: Information Security, Integrity, Encryption, Decryption, Symmetric Algorithm, Cipher, Storing Password

## I.INTRODUCTION

In cyber security, Encipher \& Decode is used to maintain data purity \& Private. Cryptography is all about encryption \& decryption. And cryptography comes under cryptology. In Cryptography, algorithms are there, through which safety takes place. Algorithms are of 2 types: 1) Symmetric Algorithm \& 2) Asymmetric algorithm. In symmetry algorithm, public key encoding concept is used \& in Asymmetric algorithm, public key - private key algorithm concept is used. According to security tester, symmetric key algorithm process fast as compare to public key (asymmetric key) algorithm because, symmetric algorithm can't use lengthy mathematical logics \& concepts. So, here I came with a new method of key algorithm to cipher a plain key to make secure communication. It's just converting a key to encrypted form. It is a method which converts 16 characters to 192 characters. Yes, it is like ciphering 128 bits to 1536 bits of encryption. And I'm using here $8 \times 8$ of matrix. It Means 64 characters is there; 26 Lowercase (small) alphabets, 26 Uppercase (CAPITAL) alphabets, $0-9$ ( 10 numerical digits), 2 special characters. And I'm translating by doing 12 rounds of matrix \& their mixture.
This is a symmetric key algorithm, where the key will be in encoded form, and only the receiver who has this application or the one who knows the algorithm flow can crack it. This algorithm is a kind of Playfair algorithm, Caesars Cipher and ROT13 algorithm but not as same as that. I referred interconnected kind of ideas of Playfair Algorithm, Caesars Cipher and ROT13 algorithm.
Password manager proceed like a digital safe, which securely supply user's usernames, passwords and other sensitive details. The usernames and passwords are encoded before saving them into the database using cryptography methods.
A password manager is mainly an encipher digital vault that supply secure password login details you use to proceed towards apps and accounts on your mobile device, websites and other services.
A password manager is a program that store all your passwords, as well as other detail, in one acceptable (suitable) location with one master password. The advantages of using a password manager are: A password manager will do the task of generating the complex passwords you need to help secure your online accounts.

## II. WORKING FLOW



Insert the data and perform the data pre-processing steps. Executes cryptography method to store the input data securely

## III. RESULT AND DISCUSSION

- Fig 1 shows that the index page in that it asking us to input data. Input types are Application Name, Email, Mobile Number, User Name, Password. And after submitting it will take to other page
- After adding inputs and submitting one message will come that is "Record Added Successfully" as shown in Fig 2.
- Fig 3 is having 3 buttons are 1) DECRYPT 2) UPDATE 3) DELETE if we click on decrypt button it show the password in decrypted form usually no one will show the password in decrypt format. And it generates key automatically which will be hidden.

Figure 1: Input Types


Figure 2: Adding Records And Storing Data in encrypted form


## VI. PROBLEM WITH PROPOSED SYSTEM

I've created this new algorithm which is built on playfair algorithm. The trouble with playfair algorithm is that, key will easily know by any cracker, because while encrypting that key will be put first row, then remaining blanks will be filled with remaining characters. Like if key is "MONARCHY". Then it will be stored in $5 * 5$ matrix is like:

| M | O | N | A | R |
| :--- | :--- | :--- | :--- | :--- |
| C | H | Y | B | D |
| E | F | G | I | K |
| L | P | Q | S | T |
| U | V | W | X | Z |

J is not there in this table but it will be with I like this I/J But in my matrix, key stored at different places like 1st, middle and last, column-wise \& rows-wise. It's not the same like Playfair, kind of, like changing positions. Cracking playfair cipher possibilities are

625 (25*25).

## V. PROCEDURE AND ALGORITHM (Encryption)

1) Take a string of set of all Uppercase(A-Z) Alphabets, Lowercase(a-z) alphabets, Numerical (0-9) digits, Special characters (@, \#) these are the string which I've taken.
e.g.: Str "ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789@\#"
2) Convert string \& store it into 1D-Array (1 Dimensional array)
3) Convert 1D-Array into 2D-array (2 Dimensional array)
4) 

| A | B | C | D | E | F | G | H |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| I | J | K | L | M | N | O | P |
| Q | R | S | T | U | V | W | X |
| Y | Z | a | b | c | d | e | f |
| g | h | i | j | k | l | m | n |
| o | p | q | r | s | t | u | v |
| w | x | y | z | 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 | 9 | 0 | $\#$ | $@$ |

Table1: Conversion of 1-Dimensional array to 2-Dimensional array
5) Take characters from particular positions $(0,3,4,7)$ row-wise \& column-wise, and store it into 2 D -Array $\mathrm{P}=$ array of position (0,3,4,7)
A=1D-array
$\mathrm{K}=\mathrm{key}$
loop I of A row-wise
loop X of P
if $\mathrm{P}=\mathrm{I}$
loop J of A [ I ] column-wise
loop Y of P

$$
\text { if } \begin{aligned}
\mathrm{J} & =\mathrm{Y} \\
\mathrm{~K} & =\mathrm{K}+\text { value at position }[\mathrm{I}][\mathrm{J}]
\end{aligned}
$$

At the end we'll get 16 digit key. And it will be stored in 2D-Array.
6) $\quad$ Swap $1^{\text {st }}$ row with $4^{\text {th }}$ Row \&
$5^{\text {th }}$ Row with $8^{\text {th }}$ Row
repeat step 4
7) Swap $1^{\text {st }}$ column with $4^{\text {th }}$ column \&
$5^{\text {th }}$ column with $8^{\text {th }}$ column
repeat step 4
8) Swap $1^{\text {st }}$ row values with last row values,
$2^{\text {nd }}$ row with (last -1 ) row,
$3^{\text {rd }}$ row with (last -2 ) row,
$4^{\text {th }}$ row with (last -3 ) row,
and repeat step 4
9) Swap $1^{\text {st }}$ column with last column,
$2^{\text {nd }}$ column with (last -1 ) column,
$3^{\text {rd }}$ column with (last -2 ) column, $4^{\text {th }}$ column with (last -3 ) column, and repeat step 4
10) Swap $1^{\text {st }}$ row values with $5^{\text {th }}$ row values,
$2^{\text {nd }}$ row with $6^{\text {th }}$ row,
$3^{\text {rd }}$ row with $7^{\text {th }}$ row,
$4^{\text {th }}$ row with $8^{\text {th }}$ row,
and repeat step 4
11) Swap $1^{\text {st }}$ column with $5^{\text {th }}$ column, $2^{\text {nd }}$ column with $6^{\text {th }}$ column,
$3^{\text {rd }}$ column with $7^{\text {th }}$ column, $4^{\text {th }}$ column with $8^{\text {th }}$ column, and repeat step 4
12) Now, in selected block, increase ASCII values with 1

| A | B | C | D | E | F | G | H |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| I | J | K | L | M | N | O | P |
| Q | R | S | T | U | V | W | X |
| Y | Z | a | b | c | d | e | f |
| g | h | i | j | k | l | m | n |
| o | p | q | r | s | t | u | v |
| w | x | y | z | 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 | 9 | 0 | $\#$ | $@$ |

Table 2: Select block from array

## And repeat step 4

13) Now, in selected block, increase ASCII values with 2

| A | B | C | D | E | F | G | H |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| I | J | K | L | M | N | O | P |
| Q | R | S | T | U | V | W | X |
| Y | Z | a | b | c | d | e | f |
| g | h | i | j | k | l | m | n |
| o | p | q | r | s | t | u | v |
| w | x | y | z | 0 | l | 2 | 3 |
| 4 | 5 | 6 | 7 | 8 | 9 | $@$ | $\#$ |

Table 3: Select block from array
And repeat step 4
14) Now, in selected block, increase ASCII values with 1


Table 4: Iincreasing Array After ASCII values
And Repeat step 4
15)

Now, in selected block, increase ASCII values with 2


Table 5: Increasing Array After ASCII values
And Repeat step 4
16) Again repeat step 14 , increase ASCII value with +2
17) After this all, we"ll get 2D-array, full of 12 keys from 12 rounds, as follows:
a. ADEHY12569aduxy\#
b. Y125ADEHuxy\#69ad
c. 1Y52DAHExu\#y96da
d. 96daxu\#yDAHE1Y52
e. ad69y\#uxEHAD25Y1
f. EHAD25Y1ad69y\#ux
g. ADEHY12569aduxy\#
h. BEEHZ22569aduxy\#
i. BEGJZ24769aduxy\#
j. BEGJZ2477:advyy\#
k. BEGJZ2477:cfvy\{\%

1. BEGJZ2477:cfvy \}"

Now, we've to joint this all keys. So, we can ge encoded keys. But, before uniting of all keys, we have to make it complex. Because last key is needed to start decoding \& $1^{\text {st }}$ one to get plain key.

Now, replace $1^{\text {st }}$ position with $6^{\text {th }}$ position $\& 7^{\text {th }}$ position with $12^{\text {th }}$ position. Will be getting following results.

1. EHAD25Y1ad69y\#ux
2. Y125ADEHuxy\#69ad
3. 1Y52DAHExu\#y96da
4. 96daxu\#yDAHE1Y52
5. ad69y\#uxEHAD25Y1
6. ADEHY12569aduxy\#
7. BEGJZ2477:cfvy ${ }^{\text {" }}$
8. BEEHZ22569aduxy\#
9. BEGJZ24769aduxy\#
10. BEGJZ2477:advyy\#
11. BEGJZ2477:cfvy\{\%
12. ADEHY12569aduxy\#

Now, unite all keys, 192-byte encrypted key will be generated, which look like:
EHAD25Y1ad69y\#uxY125ADEHuxy\#69ad1Y52DAHExu\#y96da96daxu\#yDAHE1Y52ad69y\#uxEHAD25Y1ADE HY12569aduxy\#BEGJZ2477:cfvy)"BEEHZ22569aduxy\#BEGJZ24769aduxy\#BEGJZ2477:advyy\#BEGJZ2477:cfv y\{\%ADEHY12569aduxy\#

## VI. PROCEDURE FOR ALGORITHM (Decryption)

192-byte key needed first,
EHAD25Y1ad69y\#uxY125ADEHuxy\#69ad1Y52DAHExu\#y96da96daxu\#yDAHE1Y52ad69y\#uxEHAD25Y1ADE
HY12569aduxy\#BEGJZ2477:cfvy \}"BEEHZ22569aduxy\#BEGJZ24769aduxy\#BEGJZ2477:advyy\#BEGJZ2477:cfv y\{\%ADEHY12569aduxy\#

Convert this one string into 2D-array, as follows:

1. EHAD25Y1ad69y\#ux
2. Y125ADEHuxy\#69ad
3. 1Y52DAHExu\#y96da
4. 96daxu\#yDAHE1Y52
5. ad69y\#uxEHAD25Y1
6. ADEHY12569aduxy\#
7. BEGJZ2477: cfvy \}"
8. BEEHZ22569aduxy\#
9. BEGJZ24769aduxy\#
10. BEGJZ2477:advyy\#
11. BEGJZ2477:cfvy\{\%
12. ADEHY12569aduxy\#

From 192-byte key, last 16-digit key is needed to start decryption.
But before it, we"ve to put values at their own position. Swap $1^{\text {st }}$ values with $6^{\text {th }}$ value $\& 7^{\text {th }}$ value with $12^{\text {th }}$ value.
It will look like,

1. ADEHY12569aduxy\#
2. Y125ADEHuxy\#69ad
3. 1Y52DAHExu\#y96da
4. 96daxu\#yDAHE1Y52
5. ad69y\#uxEHAD25Y1
6. EHAD25Y1ad69y\#ux
7. ADEHY12569aduxy\#
8. BEEHZ22569aduxy\#
9. BEGJZ24769aduxy\#
10. BEGJZ2477:advyy\#
11. BEGJZ2477:cfvy $\{\%$
12. BEGJZ2477:ehvy\}"

Now, pick $12^{\text {th }}$ position value, and proceed to decryption.

1) Convert key string into 1D-array.
2) Make a loop of $4 * 4$

| B | E | G | J |
| :---: | :---: | :---: | :---: |
| Z | 2 | 4 | 7 |
| 7 | $:$ | e | H |
| V | y | $\}$ | $"$ |

Table6: Array of key
3) Remove table ASCII values with -2 Key: BEGJZ2477:cfvy \{\%
4) Again, remove table ASCII values with-2 Key: BEGJZ2477:advyy\#
5) Remove table ASCII values with -1 Key: BEGJZ24769aduxy\#
6) Remove table ASCII values with - 2 Key: BEEHZ22569aduxy\#
7) Remove table ASCII values with -1 ADEHY12569aduxy\#
8) Switch left 2 columns with right 2 columns, like swapping $1^{\text {st }}$ position table-column with $3^{\text {rd }}$ position table-column $\& 2^{\text {nd }}$
position table-column with $4^{\text {th }}$ position table-column Key: EHAD25Y1ad69y\#ux
9) Swap top 2 rows with bottom 2 rows, like swapping $1^{\text {st }}$ position table-row with $3^{\text {rd }}$ position table-row $\& 2^{\text {nd }}$ position tablerow with $4^{\text {th }}$ position table-row Key: ad69y\#uxEHAD25Y1
10) Swap first column with fourth column, second column with $3^{\text {rd }}$ column Key: 96daxu\#yDAHE1Y52
11) Swap first row with fourth row, second row with third row Key: 1Y52DAHExu\#y96da
12) Swap $1^{\text {st }}$ column with $2^{\text {nd }}$ column, $3^{\text {rd }}$ column with $4^{\text {th }}$ column Key: Y125ADEHuxy\#69ad
13) Swap first row with second row, third row with fourth row

Key: ADEHY12569aduxy\#
Now finally, we got our key 16-digit key which is ADEHY12569aduxy\#

## VII. CONCLUSIONS

CipherKey creates algorithm in the strongest \& non-breakable. Plain key will be automatic generated every time. So, for decoder it's difficult to crack. By using this algorithm, probity \& clandestineness should be maintained If bits are too increased then it will be difficult to decode.

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